LOCAL FINITENESS IN VARIETIES OF MS4-ALGEBRAS

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Abstract. It is a classic result of Segerberg and Maksimova that a variety of S4-algebras is locally finite iff it is of finite depth. Since the logic MS4 (monadic S4) axiomatizes the one-variable fragment of QS4 (predicate S4), it is natural to try to generalize the Segerberg—Maksimova theorem to this setting. We obtain several results in this direction. Our positive results include the identification of the largest semisimple variety of MS4-algebras. We prove that the corresponding logic MS4_S has the finite model property. We show that both 55^2 and 54_u are proper extensions of MS4_S, and that a direct generalization of the Segerberg—Maksimova theorem holds for a family of varieties containing the variety of 54_u -algebras. Our negative results include a translation of varieties of 55_2 -algebras into varieties of MS4_S-algebras of depth 2, which preserves and reflects local finiteness. This, in particular, shows that the problem of characterizing locally finite varieties of MS4-algebras (even of MS4_S-algebras) is at least as hard as that of characterizing locally finite varieties of 55_2 -algebras—a problem that remains wide open.

§1. Introduction. It is a classic result in modal logic, known as the Segerberg–Maksimova theorem, that a normal extension of S4 is locally tabular iff it is of finite depth [25, 30]. Since a logic is locally tabular iff its corresponding variety is locally finite, this provides a characterization of locally finite varieties of S4-algebras (see Section 4).

An important extension of S4 is the bimodal logic MS4—monadic S4—which axiomatizes the one-variable fragment of predicate S4 [14]. It is known that a direct adaptation of the Segerberg–Maksimova theorem is no longer true for varieties of MS4-algebras. Indeed, the well-known variety of two-dimensional diagonal-free cylindric algebras is not locally finite [18, Theorem 2.1.11], although it is precisely the variety of MS4-algebras of depth 1. This variety corresponds to the well-known bimodal system S5², which axiomatizes the (diagonal-free) two-variable fragment of classical predicate logic (or, equivalently, the one-variable fragment of predicate S5). Things improve when the attention is restricted to the bimodal logic MGrz—the one-variable fragment of predicate Grz (Grzegorczyk's well-known modal logic [17]). In this case, the Segerberg–Maksimova theorem has a direct generalization: a variety V of MGrz-algebras is locally finite iff it is of finite depth [5, Section 4.10].

It is natural to seek a characterization of locally finite varieties of MS4-algebras. In the special case of S5²-algebras, such a characterization was given in [6, Section 4], where it was shown that every proper subvariety of the variety S5² of all S5²-algebras

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is locally finite. As we pointed out above, S5²-algebras are exactly MS4-algebras of depth 1. One of our main results (Section 6) shows that already a characterization of locally finite varieties of S5²-algebras of depth 2 is at least as hard as characterizing locally finite varieties of S5²-algebras of depth 2 is at least as hard as characterizing locally finite varieties of S5²-the variety corresponding to the fusion of S5 with itself. This problem, as well as the related problem of characterizing locally finite varieties of (unimodal) KTB-algebras, remains wide open (see [31] for a recent treatment). This we do by demonstrating a translation of subvarieties of S5₂ into subvarieties of MS4[2]—the variety of MS4-algebras of depth 2—that preserves and reflects local finiteness. Thus, already characterizing locally finite subvarieties of MS4[2] would solve the problem of local finiteness for varieties of S5₂-algebras. In addition, we show that our translation does not naturally extend to depth 3 or higher, suggesting that in the general case local finiteness in MS4 might be even more difficult. To summarize, characterizing locally finite varieties of MS4-algebras is a hard problem.

On the positive side, we show in Section 4 that the Segerberg–Maksimova theorem has an obvious generalization to a family of varieties containing $S4_u$ which corresponds to the well-known bimodal logic $S4_u$ —S4 with the universal modality—a logic that plays an important role in the study of the region-based theory of space (see, e.g., [1]). Both $S5^2$ and $S4_u$ are semisimple subvarieties of MS4. We identify in Section 3 the largest semisimple subvariety of MS4, which we denote by $MS4_S$. We prove that the corresponding bimodal logic $MS4_S$ has the finite model property. Unfortunately, it is a hard problem to characterize local finiteness already in subvarieties of $MS4_S$ as our translation lands in $MS4_S[2]$.

These difficulties manifest themselves in the structure of dual spaces of finitely generated MS4-algebras. While the well-known coloring technique [12] does generalize to MS4-algebras, unlike finitely generated S4-algebras, the dual spaces of finitely generated MS4-algebras may have infinitely many points of finite depth. We demonstrate this by describing the dual space of the well-known Erdös–Tarski algebra, which is an infinite one-generated S5²-algebra, thus has infinitely many points of depth 1. Some related open problems are posed in Section 5, and possible future directions of research are discussed in Section 7.

§2. Preliminaries. Let \mathcal{L} be a propositional modal language with two modalities \Diamond and \exists . As usual, we write $\Box = \neg \Diamond \neg$ and $\forall = \neg \exists \neg$.

DEFINITION 2.1 [14]. The bimodal logic MS4 is the smallest normal modal logic in \mathcal{L} containing

- The S4 axioms for \Diamond (i.e., the K axiom along with $p \to \Diamond p$ and $\Diamond \Diamond p \to \Diamond p$),
- The S5 axioms for \exists (e.g., the S4 axioms along with $\exists \forall p \rightarrow \forall p$),
- The left commutativity axiom $\exists \Diamond p \rightarrow \Diamond \exists p$.

Remark 2.2. In the terminology of [15], $MS4 = [S4, S5]^{EX}$; that is, MS4 is the expanding relativized product of S4 and S5.

Algebraic semantics for MS4 is given by the following Boolean algebras with operators, first considered by Fischer–Servi [14] under the name of *bimodal algebras*.

DEFINITION 2.3. An MS4-algebra is a tuple $\mathfrak{A} = (B, \lozenge, \exists)$ such that

• (B, \lozenge) is an S4-algebra, i.e., B is a Boolean algebra and \lozenge is a unary function on B satisfying the identities of a closure operator:

$$\Diamond 0 = 0 \qquad \Diamond (a \vee b) = \Diamond a \vee \Diamond b \qquad a \leq \Diamond a \qquad \Diamond \Diamond a \leq \Diamond a.$$

- (B, \exists) is an S5-algebra, that is an S4-algebra that in addition satisfies $\exists \forall a \leq \forall a$ (where we write for Boolean negation and use standard abbreviations $\Box = -\lozenge -$ and $\forall = -\exists -$).
- $\exists \Diamond a \leq \Diamond \exists a$.

It is clear that the class of MS4-algebras is equationally definable, and hence forms a variety. We denote it and the corresponding category by **MS4**. The last axiom in the preceding definition has a few equivalent forms that we will make use of in the sequel.

Lemma 2.4. *Each of the following identities is equivalent to the axiom* $\exists \lozenge \leq \lozenge \exists$:

- (1) $\exists \Diamond \exists = \Diamond \exists (\exists preserves \Diamond -fixpoints).$
- (2) $\forall \Box \forall = \Box \forall \ (\forall \ preserves \Box -fixpoints).$
- $(3) \exists \square < \square \exists$.
- (4) $\Diamond \forall \Diamond = \forall \Diamond \ (\Diamond \ preserves \ \forall \text{-fixpoints}).$

PROOF. That the original axiom follows from (1) is obvious since \exists is increasing. For the same reason, the axiom implies the (\geq) direction in (1), while the (\leq) direction follows since \exists is idempotent. By taking Boolean negation, we see that (2) is equivalent to (1). That the axiom is equivalent to (3) is known, but not immediate; see [3, Lemma 2.5]. Taking Boolean negation yields that (3) is equivalent to the inequality $\Diamond \forall \leq \forall \Diamond$, which in turn is equivalent to (4) in the same manner that (1) is equivalent to the original axiom.

DEFINITION 2.5. For an MS4-algebra $\mathfrak{A} = (B, \Diamond, \exists)$ let B_0 be the set of \exists -fixpoints.

If \Diamond_0 is the restriction of \Diamond to B_0 , then it is straightforward to see that (B_0, \Diamond_0) is an S4-subalgebra of (B, \Diamond) . We denote it by \mathfrak{A}_0 .

Write H_{\square} for the set of \square -fixpoints of an S4-algebra (B, \lozenge) . It is well known (see, e.g., [13, Proposition 2.2.4]) that H_{\square} is a bounded sublattice of B which is a Heyting algebra, with Heyting implication given by $a \to b = \square(-a \lor b)$. The Heyting algebra H_{\square} is sometimes called the *skeleton* of (B, \lozenge) , and it is well known that, up to isomorphism, each Heyting algebra arises as the skeleton of some S4-algebra (see, e.g., [13, Section 2.5]). This plays a key role in proving faithfulness of the well-known Gödel translation of intuitionistic logic into S4 (see, e.g., [29, Section XI.8]).

REMARK 2.6. Similar to monadic Heyting algebras (see, e.g., [8, Section 3]), we have that MS4-algebras can be represented as pairs of S4-algebras $(\mathfrak{A}, \mathfrak{A}_0)$ such that \mathfrak{A}_0 is an S4-subalgebra of \mathfrak{A} and the inclusion $\mathfrak{A}_0 \hookrightarrow \mathfrak{A}$ has a left adjoint (\exists) .

We will make use of Jónsson–Tarski duality [20] to work with the dual spaces of MS4-algebras. We recall some of the relevant notions:

DEFINITION 2.7. A topological space *X* is a *Stone space* if it is compact, Hausdorff, and zero-dimensional—that is, *X* has a basis of *clopen sets* (sets that are both closed and open).

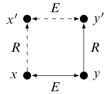
DEFINITION 2.8. Let X be a Stone space. We say that a relation $R \subseteq X^2$ is *continuous* if

- (1) $R(x) := \{ y \in X : xRy \}$ is closed for each $x \in X$ (R is point-closed).
- (2) $R^{-1}(U) := \{ y \in X : yRx \text{ for some } x \in U \}$ is clopen whenever $U \subseteq X$ is clopen.

Specializing multimodal descriptive general frames (see, e.g., [9, Section 5.5]) to MS4 yields.

DEFINITION 2.9. A descriptive MS4-frame is a tuple $\mathfrak{F} = (X, R, E)$ such that X is a Stone space, R is a quasi-order (reflexive and transitive), and E is an equivalence relation on X such that

- both R and E are continuous relations.
- $RE \subseteq ER$ —that is, $\forall x, y, y' \in X$ (xEy and yRy') $\rightarrow \exists x' \in X$ (xRx' and x'Ey').



We will refer to descriptive frames simply as *frames* because they are the only kind of frames we will deal with in this paper (as opposed to Kripke frames or the broader class of general frames studied in modal logic [9, 11]).

REMARK 2.10. The subalgebra $\mathfrak{A}_0 = (B_0, \Diamond_0)$ of \exists -fixpoints (Definition 2.5) is dually identified with the S4-frame $(X/E, \overline{R})$ of E-equivalence classes, where

$$\alpha \overline{R} \beta$$
 iff $\exists x \in \alpha, y \in \beta : xRy$.

The condition $RE \subseteq ER$ ensures that this is a well-defined quasi-order. In fact, this is a general construction arising from any *correct partition* of a descriptive frame [12]. We discuss this in more detail in Section 5.

The next definition is well known (see, e.g., [9, Section 3.3]).

DEFINITION 2.11. Let X, X' be sets, S a binary relation on X, and S' a binary relation on X'. A function $f: X \to X'$ is a p-morphism (or bounded morphism) with respect to (S, S') if

- xSy implies f(x)S'f(y);
- f(x)S'y' implies xSy for some $y \in X$ with f(y) = y'.

DEFINITION 2.12. An MS4-morphism is a continuous map $f:(X,R,E) \to (X',R',E')$ between MS4-frames such that f is a p-morphism with respect to both (R,R') and (E,E').

Specializing Jónsson–Tarski duality to MS4-algebras yields:

Theorem 2.13. The category of MS4-algebras with homomorphisms and the category of descriptive MS4-frames with MS4-morphisms are dually equivalent.

REMARK 2.14. The above dual equivalence is implemented as follows: with each MS4-frame $\mathfrak{F}=(X,R,E)$ we associate the MS4-algebra $\mathfrak{F}^*=(\operatorname{Clp} X, \Diamond_R, \Diamond_E)$ of clopen subsets of X, where \Diamond_R and \Diamond_E are the *dual operators* of R and E, defined by

$$\Diamond_R U = R^{-1}(U)$$
 and $\Diamond_E U = E(U)$.

In the other direction, with each MS4-algebra $\mathfrak{A}=(B,\Diamond,\exists)$ we associate the descriptive frame $\mathfrak{A}_*=(\mathrm{Uf}\,B,R_\Diamond,R_\exists)$ of ultrafilters of B, where R_\Diamond and R_\exists are the *dual relations* of \Diamond and \exists , defined by

$$xR_{\Diamond}y \text{ iff } x \cap H_{\square} \subseteq y \text{ and } xR_{\exists}y \text{ iff } x \cap \mathfrak{A}_0 = y \cap \mathfrak{A}_0,$$

and Uf B is equipped with the Stone topology generated by the clopen basis $\{\mathfrak{s}(a): a \in B\}$, where $\mathfrak{s}(a) = \{x: a \in x\}$. The unit isomorphisms of the duality are given by

- $\mathfrak{A} \to (\mathfrak{A}_*)^*$; $a \mapsto \{x : a \in x\}$,
- $\mathfrak{F} \to (\mathfrak{F}^*)_*; x \mapsto \{a : x \in a\}.$

The dual of a morphism f (in both categories) is the inverse image map $f^{-1}[\cdot]$, and the natural bijection $\operatorname{Hom}(\mathfrak{A},\mathfrak{F}^*)\cong\operatorname{Hom}(\mathfrak{F},\mathfrak{A}_*)$ is given by associating to $f:\mathfrak{A}\to\mathfrak{F}^*$ the morphism $\bar{f}:\mathfrak{F}\to\mathfrak{A}_*$ defined by $x\mapsto\{a\in\mathfrak{A}:x\in f(a)\}$, and to $g:\mathfrak{F}\to\mathfrak{A}_*$ the morphism $\bar{g}:\mathfrak{A}\to\mathfrak{F}^*$ defined by $a\mapsto\{x\in\mathfrak{F}:a\in g(x)\}$.

§3. Semisimple varieties. Recall [10, Section IV.12] that a variety V of algebras is semisimple if every subdirectly irreducible algebra from V is simple. It is well known that S5 is the largest semisimple subvariety of S4. In this section we introduce the largest semisimple subvariety of MS4 and its corresponding logic, which we denote by MS4_S and MS4_S, respectively. The variety MS4_S contains two well-known varieties S5² and S4_u, whose corresponding logics are S5² [15, p. 230] and S4_u [15, p. 38]. One of our main results in this section shows that MS4_S is generated by its finite algebras, and hence that MS4_S has the finite model property (and is decidable).

The standard correspondence between congruences and modal filters in BAOs (see [34, Section 4]) yields that the lattice of congruences of an MS4-algebra $\mathfrak A$ is isomorphic to the lattice of modal filters of $\mathfrak A$, where a *modal filter* of $\mathfrak A$ is a filter F satisfying $a \in F$ implies $\Box a, \forall a \in F$. Under standard duality theory (as outlined in Remark 2.14), filters of $\mathfrak A$ closed under a modal necessity operator \Box correspond dually to closed subsets of $\mathfrak A_*$ that are R_\lozenge -upsets, i.e., closed subsets U such that $x \in U$ and $xR_\lozenge y$ imply $y \in U$ (where R_\lozenge is the dual relation of \lozenge). Hence, modal filters correspond dually to closed E-saturated R-upsets (we refer to E-upsets as E-saturated sets since they are unions of equivalence classes of E). We next give a nicer characterization of congruences of $\mathfrak A$. For this we introduce an auxiliary S4-modality on $\mathfrak A$, and the corresponding auxiliary quasi-order on the dual descriptive frame of $\mathfrak A$.

Definition 3.1.

- (1) For an MS4-algebra $\mathfrak{A} = (B, \Diamond, \exists)$, define $\blacklozenge = \Diamond \exists$.
- (2) For an MS4-frame $\mathfrak{F} = (X, R, E)$, define Q = ER.

Clearly \blacklozenge is an S4-possibility operator on B and Q is a quasi-order on X. The next lemma shows that Q is the dual relation of \blacklozenge . Let $\blacksquare = \neg \blacklozenge \neg$. Then $\blacksquare = \Box \forall$ and \blacksquare is an S4-necessity operator on B. We let H_{\blacksquare} be the Heyting algebra of fixpoints of \blacksquare .

LEMMA 3.2. Let $\mathfrak A$ be an MS4-algebra and $\mathfrak F$ its dual MS4-frame. Then xQy iff $x \cap H_{\blacksquare} \subseteq y$ for each $x, y \in \mathfrak F$.

PROOF. First suppose that xQy and $a \in x \cap H_{\blacksquare}$. Then there is z such that xRz and zEy. Therefore, $x \cap H_{\square} \subseteq z$ and $z \cap B_0 = y \cap B_0$. Since $a \in H_{\blacksquare}$, we have $a = \square \forall a = \forall \square \forall a$ (where the last equality follows from Lemma 2.4(2)), so $a \in H_{\square} \cap B_0$. But then from $a \in x$ it follows that $a \in z$, which then implies that $a \in y$. Thus, $x \cap H_{\blacksquare} \subseteq y$.

Conversely, suppose that $x \cap H_{\blacksquare} \subseteq y$. Let F be the filter generated by $(x \cap H_{\square}) \cup (y \cap B_0)$ and I the ideal generated by $B_0 \setminus y$. We show that $F \cap I = \emptyset$. Otherwise there are $a \in x \cap H_{\square}$, $b \in y \cap B_0$, and $c \in B_0 \setminus y$ such that $a \wedge b \leq c$. Therefore, $a \leq b \to c$, so $a \leq \square(b \to c)$ because $a \in H_{\square}$. Since $b, c \in B_0$, we have $\square(b \to c) \in H_{\blacksquare}$. Thus, $\square(b \to c) \in x \cap H_{\blacksquare} \subseteq y$, and so $b \to c \in y$. Consequently, $c \in y$ because $b \in y$. The obtained contradiction proves that $F \cap I = \emptyset$. Therefore, there is an ultrafilter z such that $x \cap H_{\square} \subseteq z$ and $z \cap B_0 = y \cap B_0$. But then xRz and zEy, yielding that xQy.

DEFINITION 3.3. We call a filter F of an MS4-algebra a \blacksquare -filter if $a \in F$ implies $\blacksquare a \in F$.

We order the set of ■-filters by inclusion. It is then clear that it is a complete lattice (because it is closed under arbitrary intersections). The next theorem generalizes the well-known correspondence (see, e.g., [13, Sections 2.4 and 3.4]) between congruences and □-filters of S4-algebras and closed upsets of their dual S4-frames to the setting of MS4-algebras.

Theorem 3.4. Let $\mathfrak A$ be an MS4-algebra and $\mathfrak F$ its dual MS4-frame. The following complete lattices are isomorphic:

- (1) Modal filters (and congruences) of \mathfrak{A} .
- (2) \blacksquare -filters of \mathfrak{A} .
- (3) Modal filters (and congruences) of the S4-algebra \mathfrak{A}_0 .
- (4) Filters (and congruences) of the Heyting algebra H_{\blacksquare} .
- (5) Closed Q-upsets of \mathfrak{F} (under reverse inclusion).
- (6) Closed E-saturated R-upsets of \mathfrak{F} (under reverse inclusion).

PROOF. The modal filters and \blacksquare -filters are literally the same lattice. Therefore, so are the closed Q-upsets and the closed E-saturated R-upsets of \mathfrak{F} . This gives the equivalence of (1), (2) and (5), (6). Note that the \blacksquare -fixpoints of \mathfrak{A} are exactly the \square -fixpoints of \mathfrak{A}_0 , which yields the equivalence of (2) and (3), and also (4) by [13, Theorem 2.4.17]. Finally, these are equivalent to (5) by [13, Theorem 3.4.16].

Let $\mathfrak{F}=(X,R,E)$ be an MS4-frame. We call $x\in X$ a *Q-root* of \mathfrak{F} if Q(x)=X. Then \mathfrak{F} is *Q-rooted* if it has a *Q-root*, and *strongly Q-rooted* if the set of *Q-roots* is nonempty and open. A *Q-cluster* is an equivalence class of $Q\cap Q^{-1}$. The previous theorem then yields a characterization of simple and subdirectly irreducible algebras (we abbreviate subdirectly irreducible by s.i.):

THEOREM 3.5. Let \mathfrak{A} be a (non-trivial) MS4-algebra and \mathfrak{F} its dual MS4-frame.

- (1) \mathfrak{A} is s.i. iff \mathfrak{F} is strongly Q-rooted.
- (2) A is simple iff & is a Q-cluster.
- PROOF. (1) It is well known (see, e.g., [10, Section II.8]) that $\mathfrak A$ is s.i. iff $\mathfrak A$ has a least nontrivial congruence. By the above theorem, this is equivalent to $\mathfrak F$ having the largest closed Q-upset Y different from X. The complement of Y is the nonempty open set of Q-roots of $\mathfrak F$.
- (2) $\mathfrak A$ is simple iff the only nontrivial congruence is A^2 [10, Section II.8]. Thus, using (1), $\mathfrak A$ is simple iff $Y=\varnothing$, which is equivalent to $\mathfrak F$ being a Q-cluster.
- REMARK 3.6. The above theorem also follows from a general result of Venema [33]. For this it is enough to observe that Q is the reachability relation in an MS4-frame, and thus the above theorem follows from [33, Corollary 1].
- Let $\mathfrak A$ be a s.i. MS4-algebra. By the above theorem, $\mathfrak A$ is simple iff \blacklozenge is an S5-possibility operator, which happens iff $\blacklozenge \blacksquare a \leq \blacksquare a$ for each $a \in \mathfrak A$. This motivates the following definition.
- DEFINITION 3.7. Define $MS4_S = MS4 + \blacklozenge \blacksquare p \rightarrow \blacksquare p$ and its corresponding variety $MS4_S = MS4 + (\blacklozenge \blacksquare a \leq \blacksquare a)$.
 - THEOREM 3.8. **MS4**_S is the largest semisimple subvariety of **MS4**.
- PROOF. Let **V** be a semisimple subvariety of **MS4** and let $\mathfrak{A} \in \mathbf{V}$ be s.i. Then \mathfrak{A} is simple, so the dual MS4-frame \mathfrak{F} of \mathfrak{A} is a *Q*-cluster by Theorem 3.5(2). Therefore, *Q* is an equivalence relation on \mathfrak{F} , and hence $\mathfrak{A} \models \mathbf{\Phi} \blacksquare p \to \blacksquare p$. Thus, $\mathfrak{A} \in \mathbf{MS4_S}$, and so $\mathbf{V} \subseteq \mathbf{MS4_S}$.
- REMARK 3.9. The above theorem also follows from a general result of Kowalski and Kracht [21, Theorem 12 and Proposition 8] that semisimplicity of a variety of modal algebras in finite signature is equivalent to the definability of a universal modality.

The two logics $S4_u$ and $S5^2$ mentioned in the introduction are actually both extensions of $MS4_5$. We recall the definitions.

Definition 3.10.

- (1) [4, 16] S4_u is the smallest normal modal logic in \mathcal{L} containing the S4 axioms for \Diamond , the S5 axioms for \exists , and the bridge axiom $\Diamond p \to \exists p$. Algebraically, S4_u is the variety MS4 + ($\Diamond a \leq \exists a$).
- (2) [15, Section 5.1] S5² is the *product* of S5 with itself, i.e. the smallest normal modal logic in \mathcal{L} containing the S5 axioms for both \lozenge , \exists and the commutativity axiom $\lozenge \exists p \leftrightarrow \exists \lozenge p$. Algebraically, S5² is the variety **MS4** + $(\lozenge \Box a \leq \Box a)$.

The variety $S5^2$ is also known as the variety of diagonal-free cylindric algebras of dimension two [18, Definition 1.1.2]. Its subvarieties were studied in detail in [6]. As we will see in Section 4, $S5^2$ is the variety of all MS4-algebras of depth 1.

Note that $S4_u$ and $S5^2$ are incomparable. Indeed, consider the frames \mathfrak{F} and \mathfrak{G} depicted in Figure 1, where the black arrows represent R and dotted circles E. The dual algebra of \mathfrak{F} belongs to $S4_u$ since there is only one E-cluster, but evidently R is not an equivalence, so it does not belong to $S5^2$. On the other hand, the dual

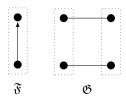


Figure 1. The dual algebra of \mathfrak{F} belongs to $S4_u$ but not to $S5^2$, and the opposite holds for the dual algebra of \mathfrak{G} .

algebra of \mathfrak{G} belongs to $S5^2$ since R and E are commuting equivalence relations, but R-clusters are not contained in E-clusters, so the identity $\Diamond a \leq \exists a$ of $S4_u$ is falsified. Consequently, both $S4_u$ and $S5^2$ are proper subvarieties of $MS4_S$.

While the characterization of locally finite subvarieties of these two varieties is known (see [6] and Corollary 4.5 below), the question is dramatically more complicated for subvarieties of MS4_S, as will be demonstrated in Section 6.

We conclude this section by establishing that $MS4_S$ is generated by its finite members, and hence that $MS4_S$ has the finite model property (hereafter abbreviated fmp). For this we recall (see, e.g., [10, Definition II.10.14]) that an algebra is *locally finite* if every finitely generated subalgebra is finite.

Lemma 3.11. If $\mathfrak{A} = (B, \Diamond, \exists)$ is an MS4_S-algebra, then $(B, \blacklozenge, \exists)$ is a locally finite S5²-algebra.

PROOF. Since 𝔄 is an MS4_S-algebra, ♦ is an S5-operator. Moreover,

$$\exists \blacklozenge a = \exists \Diamond \exists a = \Diamond \exists a = \blacklozenge \exists a$$

(the second equality follows from Lemma 2.4(1)). Therefore, $(B, \blacklozenge, \exists)$ is an S5²-algebra that in addition satisfies $\exists \blacklozenge a = \blacklozenge a$ for each $a \in B$. Thus, $(B, \blacklozenge, \exists)$ belongs to a proper subvariety of S5², which is locally finite by [6, Section 4]. Consequently, $(B, \blacklozenge, \exists)$ is locally finite.

Let $\mathfrak{A} = (B, \lozenge, \exists)$ be an MS4_S-algebra. For a finite $S \subseteq B$, let B' be the subalgebra of $(B, \blacklozenge, \exists)$ generated by S. By the above lemma, B' is finite. Let K be the fixpoints of \lozenge , define \lozenge' on B' by

$$\Diamond' a = \bigwedge \{ x \in B' \cap K : a \le x \},$$

and set $\mathfrak{A}_S = (B', \lozenge', \exists)$.

LEMMA 3.12. \mathfrak{A}_S is a finite MS4_S-algebra.

PROOF. As we already pointed out, \mathfrak{A}_S is finite by Lemma 3.11. Moreover, \lozenge' is an **S4**-operator on \mathfrak{A}_S (see [27, Lemma 2.3]). For the reader's convenience, we sketch the proof of this result.

- That $\lozenge' 0 = 0$ is clear since $0 \in B' \cap K$.
- That $\lozenge'(a \lor b) = \lozenge'a \lor \lozenge'b$ follows from distributivity since the meets involved in the definition of \lozenge' are finite.
- That $a \leq \lozenge' a$ follows from the definition of $\lozenge' a$.

• To see that $\lozenge' \lozenge' a \le \lozenge' a$, it is enough to observe that if $b \in B' \cap K$, then $a \le b$ iff $\lozenge' a \le b$.

We next show that $\exists \lozenge' a \leq \lozenge' \exists a$, or equivalently that $\exists \lozenge' \exists a = \lozenge' \exists a$, i.e. that $\lozenge' \exists a$ is a \exists -fixpoint for each $a \in \mathfrak{A}_S$. First observe that if $b \in B' \cap K$ with $\exists a \leq b$, then $\forall b \in B' \cap K$ (since $\lozenge \forall \lozenge = \forall \lozenge$) and $\exists a \leq \forall b$. Because $\forall b \leq b$, we have $\lozenge' \exists a = \bigwedge \{ \forall b : b \in B' \cap K, \exists a \leq b \}$. But this is a finite meet of \exists -fixpoints and hence an \exists -fixpoint.

Finally, we check $\blacklozenge' \blacksquare' a \leq \blacksquare' a$ for each $a \in B'$, where $\blacklozenge' = \lozenge' \exists$ and $\blacksquare' = -\blacklozenge' -$. Since $\lozenge \exists a \in B'$ (because B' is closed under \blacklozenge), $\lozenge' \exists a = \lozenge \exists a$, so $\blacklozenge' = \blacklozenge|_{B'}$, and hence the inequality holds because $\mathfrak A$ is an MS4_S-algebra.

THEOREM 3.13. MS4_S has the fmp.

PROOF. If $\mathsf{MS4}_S \not\vdash \varphi$, then there is an $\mathsf{MS4}_S$ -algebra $\mathfrak{A} = (B, \lozenge, \exists)$ and an n-tuple \overline{a} such that $\varphi(\overline{a}) \neq 1$ in \mathfrak{A} . Let S be the set of subterms of $\varphi(\overline{a})$. Since S is finite, $\mathfrak{A}_S = (B', \lozenge', \exists)$ is a finite $\mathsf{MS4}_S$ -algebra by the previous lemma. Moreover, the definition of \lozenge' ensures that if $\lozenge a \in B'$, then $\lozenge' a = \lozenge a$. Therefore, the computation of $\varphi(\overline{a})$ in \mathfrak{A}_S is identical to that in \mathfrak{A} , and hence $\varphi(\overline{a}) \neq 1$ in \mathfrak{A}_S . Thus, φ is falsified in \mathfrak{A}_S .

As a corollary we obtain the following pair of results, where the first concerns the finite-depth extensions of MS4 (a notion we will define in Section 4).

COROLLARY 3.14.

- (1) $MS4_S = \bigcap_n MS4_S[n]$.
- (2) MS4_S is decidable.

PROOF. (1) If MS4_S $\not\vdash \varphi$, then by the previous theorem, φ is falsified on a finite MS4_S-algebra \mathfrak{A} . Let n be the depth of \mathfrak{A} . Then $\mathfrak{A} \models \mathsf{MS4}_{\mathsf{S}}[n]$, so $\mathsf{MS4}_{\mathsf{S}}[n] \not\vdash \varphi$, and hence $\bigcap_n \mathsf{MS4}_{\mathsf{S}}[n] \not\vdash \varphi$.

- (2) This is obvious since $MS4_S$ is finitely axiomatizable and has the fmp.
- **§4. Local finiteness.** A classic result of Segerberg and Maksimova gives a characterization of local finiteness in the lattice of subvarieties of **S4**. By the *depth* of an S4-algebra $\mathfrak A$ or its dual frame $\mathfrak F$, we mean the longest length of a proper R-chain in $\mathfrak F$; that is, a sequence of points $x_1Rx_2R...Rx_n$ where $\neg(x_jRx_i)$ for any i < j. If there is no bound on such R-chains, then $\mathfrak A$ is of depth ω . We say a variety $\mathbf V \subseteq \mathbf S\mathbf A$ has $depth \leq n$ if the depth of each algebra in $\mathbf V$ is $\leq n$; if $\mathbf V$ contains algebras of arbitrary depth, we say it is of $depth \omega$. We will also apply these notions to (varieties of) MS4-algebras and their dual frames, understanding the depth of an MS4-algebra to be the depth of its S4-reduct (that is, the R-depth of its dual frame).

Consider the family of formulas P_n defined by

$$P_1 = \Diamond \Box q_1 \to \Box q_1 \qquad P_n = \Diamond (\Box q_n \land \neg P_{n-1}) \to \Box q_n.$$

It is well known (see, e.g., [11, Theorem 3.44]) that a variety $V \subseteq S4$ has depth $\leq n$ iff $V \models P_n$. Following [32], we write V[n] to refer to the subvariety $V + P_n$. For example, since equivalence relations are precisely quasi-orders of depth 1, we have S4[1] = S5. Similarly, $S5^2 = MS4[1]$. There is a pleasant analogy between $S5 \subseteq S4$

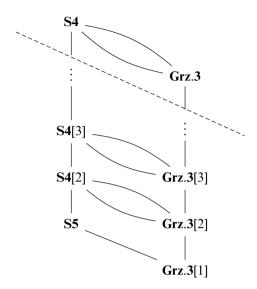


FIGURE 2. The lattice of subvarieties of S4, dually isomorphic to the lattice of normal extensions of S4. The locally finite varieties are precisely the ones below the dotted line.

and $MS4_S \subseteq MS4$ in that S5 is the largest subvariety of S4 of *R*-depth 1, while $MS4_S$ is the largest subvariety of MS4 of *Q*-depth 1. We have (see, e.g., [11, Section 12.4]):

Theorem 4.1 (Segerberg–Maksimova). A variety $V \subseteq S4$ is locally finite iff $V \models P_n$ for some $n < \omega$.

The lattice of subvarieties of S4 has the structure depicted in Figure 2. In particular, the variety Grz.3, the subvariety of S4 defined by the formulas

$$\Box(\Box(p \to \Box p) \to p) \to p \qquad \Box(\Box p \to q) \lor \Box(\Box q \to p),$$

is the least non-locally-finite subvariety of S4 (and hence local finiteness below S4 is decidable). In the figure, the varieties Grz.3[n] appearing on the right are the varieties generated by the single algebra whose dual space is the n-element chain, and Grz.3 is the variety generated by all such algebras.

It is natural to try to extend the Segerberg–Maksimova theorem to MS4. However, the theorem fails even in depth 1 since the variety $MS4[1] = S5^2$ is not locally finite: the well-known Erdös–Tarski algebra is a one-generated infinite $S5^2$ -algebra (see, e.g., [18, Theorem 2.1.11]). Thus, finite depth does not characterize local finiteness in subvarieties of MS4. To the authors' knowledge, the only positive result in this direction is that the Segerberg–Maksimova theorem holds for the variety MGrz of monadic Grzegorczyk algebras (see [5, Section 4.10]).

In this paper we will make the case that characterizing local finiteness in subvarieties of MS4 is a hard problem. In fact, we will show that a characterization of local finiteness even in subvarieties of $MS4_S[2]$ would yield one for $S5_2$. Here $S5_2$

corresponds to the fusion $S5_2 = S5 \oplus S5$, the bimodal logic of two S5 modalities with no connecting axioms (discussed in more detail in Section 6).

We do have the following criterion of local finiteness in subvarieties of **MS4**, but item (2) is in general hard to verify (as we will see shortly).

THEOREM 4.2. A subvariety V of MS4 is locally finite iff

- (1) **V** is of finite depth ($\mathbf{V} \models P_k$ for some k);
- (2) there is a function $f: \omega \to \omega$ such that for every n-generated s.i. $\mathfrak{A} \in \mathbf{V}$, the algebra \mathfrak{A}_0 is f(n)-generated as an S4-algebra.

PROOF. First suppose that **V** is locally finite. Then for every n, the free n-generated **V**-algebra is finite. Letting f(n) be its cardinality, $f:\omega\to\omega$ is our desired function, and hence (2) is satisfied. Moreover, (1) is satisfied since the **S4**-reduct of **V** is a uniformly locally finite class (that is, there is a uniform finite bound on the size of an n-generated subalgebra of an algebra from this class), and so it generates a locally finite variety (see [26, p. 285]). Thus, **V** is of finite depth by the Segerberg—Maksimova theorem.

Conversely, suppose that **V** satisfies (1) and (2). By [5, Theorem 3.7(4)], it suffices to give a finite bound on the cardinality of every n-generated s.i. $\mathfrak{A} \in \mathbf{V}$. By (2), \mathfrak{A}_0 is an f(n)-generated S4-algebra, and it is of finite depth because \mathfrak{A} is of finite depth by (1). Thus, $|B_0| \leq k(n)$ by the Segerberg–Maksimova theorem. Since B_0 is the \exists -fixpoints of \mathfrak{A} , we have that \mathfrak{A} is (n + k(n))-generated as an S4-algebra, and thus we have some m(n) such that $|B| \leq m(n + k(n))$.

We next give a few applications of this theorem where (2) can be verified. For example, one can verify (2) by induction in the **MGrz** setting, as in [5, Section. 4.10]. For another example, we need the following lemma.

Lemma 4.3. Let \mathfrak{A} be a s.i. MS4-algebra. Then, for any $k \geq 1$, \mathfrak{A} validates

$$\mathsf{alt}_k^0 = \Box \forall p_1 \vee \Box (\forall p_1 \to \forall p_2) \vee \ldots \vee \Box (\forall p_1 \wedge \forall p_2 \wedge \ldots \wedge \forall p_k \to \forall p_{k+1})$$

iff $|\mathfrak{A}_0| \leq 2^k$.

PROOF. The formula alt_k^0 is obtained from the well-known formula

$$\mathsf{alt}_k = \Box p_1 \lor \Box (p_1 \to p_2) \lor ... \lor \Box (p_1 \land p_2 \land ... \land p_k \to p_{k+1})$$

by substituting $\forall p_i$ for p_i . Let $\mathfrak{F} = (X, R, E)$ be the dual of \mathfrak{A} . Since \mathfrak{A} is s.i., \mathfrak{F} is strongly Q-rooted by Theorem 3.5(1). Because the universe of the subalgebra \mathfrak{A}_0 is exactly \forall -fixpoints of \mathfrak{A} , $\mathfrak{A} \models \mathsf{alt}_k^0$ iff $\mathfrak{A}_0 \models \mathsf{alt}_k$. By Remark 2.10, the dual space of \mathfrak{A}_0 is $\mathfrak{F}_0 = (X/E, \overline{R})$, which is a rooted descriptive S4-frame. Since alt_k is a canonical formula ([9, Definition 4.30] together with [11, Theorem 5.16]), alt_k is valid on \mathfrak{F}_0 iff it is valid on the underlying Kripke frame of \mathfrak{F}_0 ([9, Proposition 5.85]). It is well known [11, Proposition 3.45] that alt_k holds in a Kripke frame iff the number of alternatives of each point is $\leq k$ or, in the rooted transitive case, the number of points in the frame is $\leq k$. Thus, alt_k^0 holds on $\mathfrak A$ iff $|X/E| \leq k$ or, equivalently, $|\mathfrak A_0| \leq 2^k$.

THEOREM 4.4. For any $k \ge 1$ and $\mathbf{V} \subseteq \mathbf{MS4} + (\mathsf{alt}_k^0 = 1)$, \mathbf{V} is locally finite iff it is of finite depth.

PROOF. We show that in this case Theorem 4.2(2) is always satisfied and the criterion reduces to being of finite depth. Let $\mathfrak{A} \in V$ be s.i. and *n*-generated. By the above lemma, $|\mathfrak{A}_0| \leq 2^k$. Thus, \mathfrak{A}_0 is (at most) 2^k -generated.

Since $S4_u \models alt_1^0$ (by Lemma 4.3), this also gives the following generalization of the Segerberg–Maksimova theorem to varieties of $S4_u$ -algebras:

COROLLARY 4.5. $V \subseteq S4_u$ is locally finite iff it is of finite depth.

REMARK 4.6. The formulas E_i^k given in [6, Theorem 4.2] bound the number of E_i -classes in a s.i. $S5^2$ -algebra (note that $S5^2$ is semisimple, so being s.i. amounts to being simple) by k or, equivalently, the size of the subalgebra of \exists_i -fixpoints by 2^k . Therefore, the local finiteness of the subvarieties of $S5^2$ defined by the E_i^k as given in [6, Lemma 4.4] also follows from our criterion.

In the general case, however, verifying Theorem 4.2(2) is more complex. In Section 6 we demonstrate that completely characterizing local finiteness even in $MS4_S[2]$ would also do so for $S5_2$, which is an open problem as mentioned in the introduction. We do this by demonstrating a translation from subvarieties of $S5_2$ to subvarieties of $MS4_S[2]$ that preserves and reflects local finiteness. Before that, we will focus our discussion on the behavior of finitely generated MS4-algebras.

§5. Finitely generated algebras. In this section we obtain several results about finitely generated MS4-algebras, including the description of the dual space of the Erdös–Tarski algebra. We will see that they are significantly more complex than finitely generated S4-algebras. We conclude the section with several open problems about finitely generated MS4-algebras.

We start by recalling the notion of a (quasi-)maximal point (see [13, Definition 1.4.9]).

DEFINITION 5.1. Let X be a set and R a quasi-order on X. Given $A \subseteq X$, we write

$$\max A = \{x \in A : xRv, v \in A \to vRx\}$$

(so for us, max means "quasi-maximum"). We define the *n-th layer* of *X* inductively as

$$D_1 = \max X$$
 $D_{n+1} = \max \left(X - \bigcup_{i=1}^n D_i\right).$

We may speak of the layers of an S4-frame or MS4-frame, always referring to the R-layers. In the case of arbitrary quasi-orders, some or all D_i may be empty (e.g., when X contains infinite ascending chains). However, we have the following (see [13, Corollary 3.2.2 and Theorem 3.2.3]).

THEOREM 5.2. Let (X, R) be a descriptive S4-frame. Then $D_1 = \max X$ is nonempty and closed.

Clearly this holds for descriptive MS4-frames as well. We concentrate on finitely generated descriptive frames; that is, frames \mathfrak{F} whose dual \mathfrak{F}^* is a finitely generated algebra. Finitely generated descriptive S4-frames have a relatively well-understood, but complicated structure. They consist of finite "discrete" layers, possibly with limit points of infinite depth.

THEOREM 5.3 (see, e.g., [11, Section 8.6]). Let (X, R) be a finitely generated descriptive S4-frame. Then every D_n $(n < \omega)$ is finite and consists of isolated points, hence is clopen.

REMARK 5.4. The same result also holds for Heyting algebras and their dual spaces; see [7, Theorem 3.1.8].

This result does not carry over to the monadic setting. In fact, the Erdös–Tarski algebra is an infinite one-generated S5²-algebra (see [18, Theorem 2.1.11]). Thus, its dual is an infinite one-generated descriptive MS4-frame of depth 1. It follows that finitely generated descriptive MS4-frames may have limit points already in depth 1. One of the main goals of this section is to describe the dual frame of the Erdös–Tarski algebra. For this purpose, we adjust the notion of a correct partition to the setting of MS4-frames.

DEFINITION 5.5 [12]. Let $\mathfrak{F} = (X, R, E)$ be an MS4-frame. A *correct partition* of \mathfrak{F} is an equivalence relation $K \subseteq X^2$ such that

- (1) *K* is *separated* (or *Boolean*, as in [22, Definition 8.3]), i.e., any x, y with $\neg(xKy)$ can be separated by a *K*-saturated clopen set.
- (2) $RK \subseteq KR$ (K is correct with respect to R);
- (3) $EK \subseteq KE$ or, equivalently, EK = KE (K is correct with respect to E).

The quotient \mathfrak{F}/K is defined to be $(X/K, \overline{R}, \overline{E})$, where $\alpha \overline{R}\beta$ iff $\exists x \in \alpha, y \in \beta : xRy$ (or, equivalently, $\forall x \in \alpha, \exists y \in \beta : xRy$), and \overline{E} is defined similarly.

Correct partitions are the kernels of continuous p-morphisms. Hence, by standard duality theory (as outlined in Remark 2.14), the lattice of correct partitions of \mathfrak{F} ordered by inclusion is dually isomorphic to the lattice of subalgebras of \mathfrak{F}^* .

REMARK 5.6. As noted in the previous definition, when a partition K is correct with respect to the equivalence relation E, the relations actually commute and we have EK = KE. Thus, not only does E give rise to an equivalence relation \overline{E} on K-classes, but K also gives rise to an equivalence relation on E-classes, which we denote by \overline{K} . In particular, two E-classes are \overline{K} -related iff every member of each class is K-related to some member of the other. This idea is put to use in Theorem 5.13 and Lemma 6.12.

We now adapt to our setting a dual characterization of finitely generated algebras; this criterion was given in the setting of S4-algebras and Heyting algebras in [12], with an identical proof.

DEFINITION 5.7. Let $\mathfrak{F} = (X, R, E)$ be an MS4-frame, and $\overline{g} = (g_1, \dots, g_n)$ a tuple of elements of \mathfrak{F}^* . Define

(1) the coloring $c_{\bar{g}}: X \to 2^n$ by

$$c_{\bar{g}}(x)(i) = \begin{cases} 1 & x \in g_i, \\ 0 & \text{otherwise;} \end{cases}$$

(2) the relation $C_{\bar{g}} \subseteq X^2$ as the kernel of $c_{\bar{g}}$, i.e., $xC_{\bar{g}}y$ iff $c_{\bar{g}}(x) = c_{\bar{g}}(y)$. We call $c_{\bar{g}}$ the *coloring* and $C_{\bar{g}}$ the *partition induced by* \bar{g} . We will commonly drop the subscript \bar{g} when unambiguous.

We call an equivalence relation K on X non-trivial if it relates two distinct elements of X (i.e., it is strictly above the diagonal in the lattice of equivalence relations).

THEOREM 5.8 (Coloring theorem). Let \mathfrak{F} be a descriptive MS4-frame, and $\overline{g} = (g_1, \dots, g_n)$ a tuple of elements of \mathfrak{F}^* . Then \mathfrak{F}^* is not generated by \overline{g} iff there exists a non-trivial correct partition K of \mathfrak{F} with $K \subseteq C$ (that is, each class of K contains only points of the same color).

PROOF. Suppose \mathfrak{F}^* is generated by \bar{g} , and let K be a non-trivial correct partition. Then K corresponds to a proper subalgebra of \mathfrak{F}^* that couldn't contain all of the g_i . That is, there is some g_i that is not K-saturated. So there are $x \in g_i$ and $y \notin g_i$ that are K-related. Thus, xKy but $c(x) \neq c(y)$, and we've shown that any non-trivial correct partition K cannot satisfy $K \subseteq C$.

Now suppose \mathfrak{F}^* is not generated by \bar{g} . Then \bar{g} generates a proper subalgebra \mathfrak{A} , which gives rise to a non-trivial correct partition K, defined by

$$xKy$$
 iff for all $a \in \mathfrak{A}$ $(x \in a \leftrightarrow y \in a)$.

Now, if $c(x) \neq c(y)$, then $x \in g_i$ and $y \in -g_i$ for some $g_i \in \mathfrak{A}$, so $\neg(xKy)$. Thus, $K \subseteq C$.

REMARK 5.9. The previous definition and theorem hold more generally (with the same proof) for any boolean algebra with unary operators or, dually, any descriptive frame with binary relations.

With all the tools we need, we are now ready to give an explicit description of the MS4-frame that is the dual frame of the Erdös–Tarski algebra.

EXAMPLE 5.10. Let
$$\mathfrak{F} = (X, E_1, E_2)$$
 where

$$X = (\omega + 1 \times \omega + 1) - \{(\omega, \omega)\} \cup \{k_1, k_2\}.$$

The equivalence classes of E_1 are rows, the equivalence classes of E_2 are columns, and $k_1, k_2 \in E_1(0, \omega) \cap E_2(\omega, 0)$ (so that $\{k_1, k_2\}$ is the only non-trivial equivalence class of $E_1 \cap E_2$); see Figure 3.

In the remainder of the section, we demonstrate that \mathfrak{F} has the desired properties. To define the topology on X, we will describe it as a two-point compactification of the space $X_0 = (\omega + 1 \times \omega + 1) - \{(\omega, \omega)\}$, where $\omega + 1$ has the usual interval topology, and X_0 has its usual topology of a subspace of the product. As an open subset of a Stone space, X_0 is zero-dimensional, locally compact, and Hausdorff (that is, X_0 is locally Stone). Define

$$g_1 = \{(i, j) : i \le j \le \omega\} \cap X_0$$
 $g_2 = X_0 - g_1$.

Note that g_1 and g_2 are disjoint clopen sets that are neither compact nor co-compact (i.e., having a compact complement) in X_0 . In the terminology of [24, Definition 2.2], $\{g_1, g_2\}$ is a 2-star in X_0 —a pair of disjoint open sets with the property that the complement of their union is compact but the complement of each set itself is not. In [24, Theorem 2.1] it is shown that this datum gives rise to a 2-point compactification of X_0 ; specifically, we have the following basis for a compact Hausdorff topology on X:

$$\mathcal{B} = \operatorname{Clp}(X_0) \cup \{U \cup \{k_i\} : U \in \operatorname{Clp}(X_0) \text{ and } g_i - U \text{ is compact in } X_0\}.$$

(Note that \mathcal{B} is not a clopen basis.)

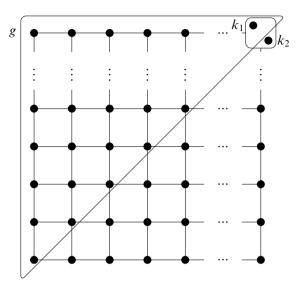


FIGURE 3. Descriptive frame of Example 5.10.

LEMMA 5.11. *X* is a Stone space.

PROOF. By [24, Theorem 2.1], X is compact Hausdorff. We show that X is zero-dimensional. For this it suffices to separate distinct points of X by a clopen subset of X (see, e.g., [19, p. 69]). First suppose that $x, y \in X_0$. Since X_0 is zero-dimensional, there is a clopen set $U \subseteq X_0$ containing x and missing y. Because X_0 is locally compact, x has a compact neighborhood K in X_0 . Therefore, there is a clopen set V with $x \in V \subseteq U \cap K$. But then V is a compact clopen of X_0 , and hence a clopen subset of X separating x and y. The points k_1 and k_2 are separated by $g_1 \cup \{k_1\}$ which is clopen of X by definition. Finally, $x \in X_0$ and k_i are separated by any compact clopen neighborhood K of x in X_0 , which is a clopen set of X.

Lemma 5.12. \mathfrak{F} is a descriptive $S5^2$ -frame.

PROOF. Since X is a Stone space by the above lemma, it remains to show that E_1, E_2 are continuous. We show that E_1 is continuous. The proof for E_2 is symmetric. Observe that all finite-index rows (i.e., $[0, \omega] \times \{j\}$ for $j < \omega$) are evidently clopen, while the "top row" $\{(i, \omega) : i < \omega\} \cup \{k_1, k_2\}$ is closed but not open (its complement is the union of all finite-index rows). Hence, E_1 is point-closed.

It is left to show that U clopen implies $E_1(U)$ is clopen. By [13, Theorem 3.1.2(IV)], it is enough to verify the following two conditions:

- (1) If $y \notin E_1(x)$, then there is a partition into two E_1 -saturated open sets separating x and y.
- (2) For all x and open V, if $E_1(x) \cap V \neq \emptyset$, then there is an open neighborhood U of x with $E_1(y) \cap V \neq \emptyset$ for all $y \in U$.

For (1), since $y \notin E_1(x)$, x and y are in distinct rows. Therefore, one of them must be in a row of finite index, say $R_n = [0, \omega] \times \{n\}$ with $n < \omega$ (if both are in rows of

finite index, take n to be the least index). Then $[0, \omega]$ and its complement provide a partition of X into two E_1 -saturated open sets separating x and y.

For (2), we consider cases: If x = (i, j) for $i, j < \omega$, $\{x\}$ is an appropriate neighborhood. If $x = (\omega, j)$ for $j < \omega$, then $[0, \omega] \times \{j\}$ is an appropriate neighborhood. Suppose $x = (i, \omega)$ for $i < \omega$. If $E_1(x) \cap V$ contains (i', ω) for $i' < \omega$, then V contains a neighborhood $\{i'\} \times [n, \omega]$ for $n < \omega$, so $\{i\} \times [n, \omega]$ is an appropriate neighborhood of x. If $E_1(x) \cap V$ contains k_1 , then V contains some $V' \cup \{k_1\}$ with $g_1 - V'$ compact in X_0 , as sets of this form are the only basic open neighborhoods of k_1 . Then there must be some n with $(g_1 - V') \cap [n, \omega]^2 = \emptyset$ (indeed any $U \subseteq X_0$ compact has this property since otherwise $\{X_0 - [n, \omega]^2 : n < \omega\}$ would be an open cover of U without a finite subcover). Therefore, $g_1 \cap [n, \omega]^2 \subseteq V' \subseteq V$, and again $\{i\} \times [n,\omega]$ is an appropriate neighborhood of x. If $E_1(x) \cap V$ contains k_2 , then by the same argument we have an *n* with $g_2 \cap [n, \omega]^2 \subseteq V$ and the same neighborhood works. Finally, suppose $x = k_s$ for $s \in \{1, 2\}$. If $E_1(x) \cap V$ contains (i, ω) for $i < \omega$, then V contains $\{i\} \times [n, \omega]$ for some n, and $(g_s \cap [n, \omega]^2) \cup \{k_s\}$ is an appropriate basic open neighborhood of x. If $E_1(x) \cap V$ contains k_t for $t \neq s$, then by the same argument as before we have $g_t \cap [n, \omega]^2 \subseteq V$ for some n, and the same neighborhood $(g_s \cap [n,\omega]^2) \cup \{k_s\}$ works.

Theorem 5.13. \mathfrak{F} is one-generated and \mathfrak{F}^* is isomorphic to the Erdös–Tarski algebra.

PROOF. Let $g = g_1 \cup \{k_1\}$ (see Figure 3). We claim $\mathfrak{F}^* = \langle g \rangle$. Suppose not. By the coloring theorem, there exists a non-trivial correct partition K of X such that every class of K is contained in g or -g. Since it is non-trivial, K must identify two points in distinct rows or columns (by assumption we cannot have k_1Kk_2). By Remark 5.6, this is to say that two distinct rows or columns must be \overline{K} -related. Let $C_i = \{i\} \times [0, \omega]$ be the i-th column and $R_j = [0, \omega] \times \{j\}$ the j-th row. Observe that C_0 cannot be related to C_i for i > 0 since otherwise a point in $-g \cap C_i$ would be K-related to a point in $-g \cap C_0 = \emptyset$.

Now suppose C_{i_1} and C_{i_2} are \overline{K} -related for $0 < i_1 < i_2 \le \omega$. Then there must be some $j_2 \in [i_1, i_2)$ so that $(i_2, j_2)K(i_1, j_1)$ for some $j_1 \in [0, i_1)$ (since all points in $-g \cap C_{i_2}$ must be K-related to points in $-g \cap C_{i_1}$). Therefore, rows j_1 and j_2 are \overline{K} -related, with $0 \le j_1 < j_2 < i_2$. Thus, $(\underline{j_2}, j_2)K(i_0, j_1)$ for some $i_0 \in [0, j_1]$. We now conclude that columns i_0 and j_2 are \overline{K} -related, and note that $0 \le i_0 < i_1$, $0 < j_2 < i_2$, and $i_0 < j_2$. So from our assumption of two \overline{K} -related columns, we have found two \overline{K} -related columns of strictly lower index, and iterating this process will eventually reveal that column 0 must be \overline{K} -related to some non-zero column, which is a contradiction by the previous paragraph.

Finally, suppose R_{j_1} and R_{j_2} are \overline{K} -related for $0 \le j_1 < j_2 \le \omega$. Then $(j_2, j_2)K(i, j_1)$ (or $k_1K(i, j_1)$ in the case $j_2 = \omega$) for some $i \in [0, j_1]$. But then columns i and j_2 are \overline{K} -related, with $i < j_2$, which is a contradiction by the last paragraph.

Consequently, such a K cannot exist, and we conclude that \mathfrak{F}^* is generated by g. There is an evident homomorphism $f:\mathfrak{F}^*\to \mathcal{P}(\omega\times\omega)$ given by $U\mapsto U\cap(\omega\times\omega)$. Then $f(g)=\{(i,j):i\leq j\}$, and f is an isomorphism onto its image (because $\omega\times\omega$ is dense in X). Thus, \mathfrak{F}^* is isomorphic to the subalgebra of $\mathcal{P}(\omega\times\omega)$ generated by f(g), which is exactly the Erdös–Tarski algebra.

This provides an example of a descriptive MS4[1]-frame whose (only) layer is infinite and contains limit points. Thus, the only part of Theorem 5.3 that has a hope to be salvaged in the monadic setting is that each layer may be clopen.

QUESTION 5.14. Let (X, R, E) be a finitely generated descriptive MS4-frame. Is $D_1 = \max X$ or any D_n a clopen (admissible, definable) set?

QUESTION 5.15. Does there exist a finitely generated descriptive MS4-frame whose finite layers consist only of limit points?

§6. Translating S5₂ to MS4_S[2]. In this section we follow up on the promise made at the end of Section 4 and give a semantic translation from subvarieties of S5₂ to subvarieties of MS4_S[2] that preserves and reflects local finiteness. Before doing so, we develop some basic facts about the variety of S5₂-algebras and their dual frames.

DEFINITION 6.1 [15, Section 3.1]. The logic S5₂ is the *fusion* S5 \oplus S5, i.e., the smallest normal modal logic in the language with two modalities \exists_1 and \exists_2 containing the S5 axioms for each \exists_i (and no other axioms).

Algebraic models of S5₂ are triples $(B, \exists_1, \exists_2)$, where (B, \exists_i) is an S5-algebra for i = 1, 2. We call such algebras S5₂-algebras. Jónsson–Tarski duality specializes to yield that the corresponding variety S5₂ is dually equivalent to the following category of descriptive frames.

DEFINITION 6.2. A descriptive S5₂-frame is a triple $\mathfrak{F} = (X, E_1, E_2)$ where X is a Stone space and E_1 , E_2 are two continuous equivalence relations on X (see Definition 2.8).

An S5₂-morphism between descriptive S5₂-frames is a continuous map $f:(X,E_1,E_2) \to (X',E_1',E_2')$ that is a p-morphism with respect to both (E_1,E_1') and (E_2,E_2') . When we depict S5₂-frames, we will depict E_1 -equivalence classes as horizontal lines and E_2 -equivalence classes with vertical analogues.

DEFINITION 6.3.

- (1) For an S5₂-algebra $(B, \exists_1, \exists_2)$, define $\diamondsuit = \exists_1 \lor \exists_2$.
- (2) For an S5₂-frame (X, E_1, E_2) , define $S = E_1 \cup E_2$.

It is straightforward to see that S is a reflexive and symmetric relation (or *edge relation*) on X and that S is the dual relation of \diamondsuit . Hence, \diamondsuit is a possibility operator validating the well-known T and B axioms, making it a KTB-modality. We write $S^* = \bigcup_{n=0}^{\infty} S^n$ for the reflexive transitive closure of S.

Definition 6.4 [33]. Let $\mathfrak{F} = (X, E_1, E_2)$ be a descriptive S5₂-frame.

- (1) Call $x \in X$ a topo-root of \mathfrak{F} if $S^*(x)$ is dense in X.
- (2) We say that \mathfrak{F} is *topo-rooted* if it has a topo-root, and *strongly topo-rooted* if it has an open set of topo-roots.

The following is a consequence of general results in [33].

THEOREM 6.5. Let \mathfrak{A} be an $S5_2$ -algebra.

- (1) \mathfrak{A} is s.i. iff \mathfrak{A}_* is strongly topo-rooted.
- (2) \mathfrak{A} is simple iff every point of \mathfrak{A}_* is a topo-root.

Following [28] (see also [23, p. 74]), we call a variety $\mathbf{V} \subseteq \mathbf{S5}_2$ *n-transitive* if $\mathbf{V} \models \diamondsuit^{n+1}a \leq \diamondsuit^n a$, and *weakly transitive* if it is *n*-transitive for some $n < \omega$. The following is a consequence of a more general result of Kowalski and Kracht [21, Theorem 12] (note that $\mathbf{S5}_2$ is automatically *cyclic* since each basic modality is S5, see [21, Proposition 6]).

Theorem 6.6. A variety $V \subseteq S5_2$ is semisimple iff it is weakly transitive.

It is a consequence of [33, Corollary 1] (where weak transitivity goes by the name of ω -transitivity) that in a weakly transitive variety of S5₂-algebras, roots and toporoots coincide (where x is a root of \mathfrak{F} if $S^*(x) = X$); together with the previous theorem, this yields:

Theorem 6.7. Let $V \subseteq S5_2$ be a weakly transitive variety. The following are equivalent:

- (1) $\mathfrak{A} \in \mathbf{V}$ is s.i.
- (2) $\mathfrak{A} \in \mathbf{V}$ is simple.
- (3) Every point of \mathfrak{A}_* is a root.

We are now ready to describe our translation. We start by a construction producing a descriptive MS4₅[2]-frame from a descriptive S5₂-frame.

Construction 6.8. Let $\mathfrak{F} = (X, E_1, E_2)$ be a descriptive $S5_2$ -frame. We let $\mathcal{L} = X/E_2$ be the quotient space and $\pi_{\mathcal{L}} : X \to \mathcal{L}$ the quotient map. We set $T(\mathfrak{F}) = (X \cup \mathcal{L}, R, E)$, where

- $X \cup \mathcal{L}$ is the (disjoint) union of X and \mathcal{L} ,
- E is the smallest equivalence on $X \cup \mathcal{L}$ containing E_2 along with

$$\{(x,\alpha):x\in X,\alpha\in\mathcal{L},x\in\alpha\},$$

• $R = E_1 \cup (X \times \mathcal{L}) \cup (\mathcal{L} \times \mathcal{L}).$

In the case that $|\mathfrak{F}|$ is finite, it follows that $|T(\mathfrak{F})| \leq 2|\mathfrak{F}|$ and $|T(\mathfrak{F})^*| \leq |\mathfrak{F}^*|^2$.

This construction is depicted in Figure 4. Under the definition of R, \mathcal{L} becomes an R-cluster, which we think of as a "top rail". To prove that $T(\mathfrak{F})$ is a descriptive MS4_S[2]-frame, we need the following.

LEMMA 6.9. $\pi_{\mathcal{L}}$ is a clopen map, meaning that U clopen implies $\pi_{\mathcal{L}}[U]$ is clopen.

PROOF. Suppose $U \subseteq X$ is clopen. Then $\pi_{\mathcal{L}}^{-1}(\pi_{\mathcal{L}}[U]) = E_2(U)$, which is clopen by continuity of E_2 . Since $\pi_{\mathcal{L}}$ is a quotient map, $\pi_{\mathcal{L}}[U]$ is clopen.

LEMMA 6.10. $T(\mathfrak{F})$ is a descriptive MS4_S[2]-frame; furthermore, $T(\mathfrak{F})$ is *Q*-rooted.

PROOF. Note that E_2 is a separated partition of X (but in general not correct with respect to E_1 in the sense of Definition 5.5) since any $(x, y) \notin E_2$ can be separated by a E_2 -saturated clopen (this is a property of any continuous quasi-order; see [13, Theorem 3.1.2]). This ensures \mathcal{L} is a Stone space [22, Lemma 8.4]. Thus, so is $X \cup \mathcal{L}$ (the union is disjoint).

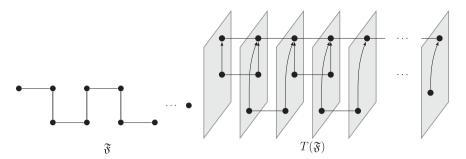


FIGURE 4. Constructing an MS4_S[2]-frame from an S5₂-frame. In the S5₂-frame on the left, E_1 -clusters are horizontal lines while E_2 -clusters are bold vertical lines. In the $MS4_S[2]$ -frame on the right, R-clusters are horizontal lines, proper R-arrows are drawn with arrowheads, and E-clusters are given by the shaded rectangles.

Clopen sets in $T(\mathfrak{F})$ are of the form $U \cup V$, with U clopen in X and V clopen in L. We have

- $E(x) = E_2(x) \cup \{E_2(x)\} \text{ for } x \in X;$
- $E(\alpha) = \alpha \cup \{\alpha\}$ for $\alpha \in \mathcal{L}$; $E(U \cup V) = E_2(U) \cup \pi_{\mathcal{L}}[U] \cup V \cup \pi_{\mathcal{L}}^{-1}(V)$.

The first two items show that E-clusters are closed. In the third item, every set on the right hand side is clopen, by continuity of E_2 and the fact that $\pi_{\mathcal{L}}$ is a clopen map (Lemma 6.9). Thus, E is continuous. For R, we have

- $R(x) = E_1(x) \cup \mathcal{L}$ for $x \in X$;
- $R(\alpha) = \mathcal{L}$ for $\alpha \in \mathcal{L}$;
- $R^{-1}(U \cup V) = \begin{cases} E_1(U) & V = \varnothing, \\ X \cup \mathcal{L} & \text{otherwise.} \end{cases}$

By the first two items, R is point-closed, which together with the third yields that R is continuous.

Now observe that Q = ER is the total relation on $T(\mathfrak{F})$ (i.e., $Q = (X \cup \mathcal{L})^2$) because for each $x \in X$ we have $\mathcal{L} \subseteq R(x)$, so $X \cup \mathcal{L} = ER(x)$. Therefore, every point of $T(\mathfrak{F})$ is a Q-root and, in particular, $RE \subseteq ER$. Finally, $T(\mathfrak{F})$ is of R-depth 2 by construction. Thus, $T(\mathfrak{F})$ is a Q-rooted descriptive MS4_S[2]-frame.

By construction, $T(\mathfrak{F})$ has exactly two layers $D_1 = \mathcal{L}$ and $D_2 = X$. By design, we have $E|_{D_2} = E_2$ and $R|_{D_2} = E_1$. Since D_1, D_2 are clopen, we can consider the relativization

$$\mathfrak{A}_i = (\operatorname{Clp} D_i, \cap, \cup, (-)', \varnothing, D_i, \lozenge', \exists'),$$

where the relative complementation and modal operators are given by

$$a' = D_i - a$$
 $\Diamond' a = \Diamond a \cap D_i$ $\exists' a = \exists a \cap D_i$.

The maps $a \mapsto a \cap D_i$ are Boolean homomorphisms. Considering the restrictions of R and E to each layer, we see that \mathfrak{A}_1 is an S5²-algebra (and in fact belongs

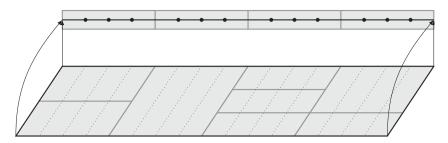


FIGURE 5. Lifting a correct partition of \mathfrak{F} to one of $T(\mathfrak{F})$, as in Lemma 6.12: *E*-classes are depicted by dotted lines (as E_2 -classes of \mathfrak{F}) with the E_2 -class itself depicted as a solid dot above on the top rail. The correct partition \widehat{K} is depicted by the shaded rectangles.

to a proper and hence locally finite subvariety of $S5^2$), while \mathfrak{A}_2 is an $S5_2$ -algebra isomorphic to \mathfrak{F}^* :

LEMMA 6.11. $\mathfrak{A}_2 \cong \mathfrak{F}^*$, and the relative operations of \mathfrak{A}_2 are definable in $T(\mathfrak{F})^*$.

PROOF. This follows from the explicit description of R^{-1} and E on clopen sets in Lemma 6.10. Also, the relative complementation and modal operators are definable using the clopen set $D_2 \in T(\mathfrak{F})^*$, e.g., by using precisely the definitions from the preceding paragraph.

Recall from Remark 5.6 that, since E_2 is an equivalence, a correct partition K on \mathfrak{F} yields an equivalence relation \overline{K} on the set \mathcal{L} of E_2 -classes of \mathfrak{F} .

Lemma 6.12. Let K be a correct partition of \mathfrak{F} . Define $\widehat{K} = K \cup \overline{K}$, a relation on $T(\mathfrak{F})$. That is, \widehat{K} is the relation K on $D_2 = X$, and the relation \overline{K} on $D_1 = \mathcal{L}$ (and relates no points on different layers); see Figure 5.

- (1) \widehat{K} is a correct partition of $T(\mathfrak{F})$.
- (2) $T(\mathfrak{F}/K) \cong T(\mathfrak{F})/\widehat{K}$.

PROOF. (1) We first show \widehat{K} is separated. If $x, y \in D_2$ and $\neg(x\widehat{K}y)$, then $\neg(xKy)$ and they can be separated by a K-saturated clopen subset of D_2 , which is also \widehat{K} -saturated (since \widehat{K} relates no points in different layers). Any $x \in D_2$ and $\alpha \in D_1$ can be separated by D_2 which is \widehat{K} -saturated by the same reason. Now suppose $\alpha, \beta \in D_1$ and $\neg(\alpha\widehat{K}\beta)$. Then $\neg(xKy)$ for any $x \in \alpha, y \in \beta$, and so any such x, y can be separated by a K-saturated clopen $A \subseteq D_2$. By Lemma 6.9, $\pi_{\mathcal{L}}[A]$ is a clopen subset of D_1 separating α and β . To see it is \widehat{K} -saturated, let $\gamma \in \pi_{\mathcal{L}}[A]$ and $\gamma \widehat{K}\delta$. Then there is $x \in A \cap \gamma$ and $y \in \delta$ with xKy (x may be chosen from x by the remarks at the end of Definition 5.5). Since x is $x \in A \cap \delta$, and hence $x \in \mathcal{L}[A]$.

We now show \widehat{K} is correct with respect to R. Suppose $x, y \in D_2$. If $x\widehat{K}yRz$ for $z \in D_2$, then $xKyE_1z$, so $xE_1y'Kz$ by correctness of K, and hence $xRy'\widehat{K}z$. If $x\widehat{K}yR\alpha$ for $\alpha \in D_1$, then $xR\alpha\widehat{K}\alpha$ (by definition of R). Suppose $\alpha, \beta \in D_1$. If

 $\alpha \widehat{K} \beta R \gamma$ for $\gamma \in D_1$ (the only possibility), then $\alpha R \gamma \widehat{K} \gamma$ (again by definition of R). In any case, $R \widehat{K} \subseteq \widehat{K} R$.

Finally, we show \widehat{K} is correct with respect to E. Suppose $x, y \in D_2$. If $x\widehat{K}yEz$ for $z \in D_2$, then $xKyE_2z$, so $xE_2y'Kz$, and hence $xEy'\widehat{K}z$. If $x\widehat{K}yE\alpha$ for $\alpha \in D_1$, then $\alpha = E_2(y)$, and $E_2(x)\overline{K}E_2(y)$. So $x \in E_2(x)\widehat{K}\alpha$. Suppose $\alpha, \beta \in D_1$. If $\alpha\widehat{K}\beta Ez$ for $z \in D_2$ (the only non-trivial possibility), then $\beta = E_2(z)$. So we must have zKx for some $x \in \alpha$, and hence $\alpha Ex\widehat{K}z$. In any case, $E\widehat{K} \subseteq \widehat{K}E$. We conclude that \widehat{K} is a correct partition of $T(\mathfrak{F})$.

(2) We exhibit an isomorphism $f: T(\mathfrak{F}/K) \to T(\mathfrak{F})/\widehat{K}$. Let f(K(x)) = K(x) on K-classes (i.e., the identity on the common second layer of both frames) and $f(\overline{E_2}(K(x))) = \overline{K}(E_2(x))$ on the additional points. To check continuity, it suffices to check on clopen subsets of the top layer only (since \widehat{K} identifies no points across layers and f is the identity on the bottom layers); such a clopen can be identified with a \overline{K} -saturated clopen U of \mathfrak{F} , and we have

$$f^{-1}(U) = \{\overline{E_2}(K(x)) : E_2(x) \in U\}$$
 is clopen in $T(\mathfrak{F}/K)$ iff $\{K(x) : E_2(x) \in U\}$ is clopen in \mathfrak{F}/K iff $\{x : E_2(x) \in U\}$ is clopen in \mathfrak{F} .

The last set is precisely $\pi_{\mathcal{L}}^{-1}(U)$, which is clopen since U is. Clearly f is surjective; for injectivity, suppose $\overline{E_2}(K(x)) \neq \overline{E_2}(K(y))$. Then

$$\neg (K(x) \overline{E_2} K(y)) \rightarrow \neg (xKy) \rightarrow \neg (E_2(x) \overline{K} E_2(y)),$$

and hence $\overline{K}(E_2(x)) \neq \overline{K}(E_2(y))$.

Since f is a continuous bijection, it remains to show it preserves and reflects the relations [13, Proposition 1.4.15]. For R, note (checking only the cases where f is not the identity on some input):

- $K(x) R E_2(K(y)) \leftrightarrow \widehat{K}(x) R \widehat{K}(E_2(y))$ (both are always true by definition).
- $E_2(K(x))$ R $E_2(K(y)) \leftrightarrow \widehat{K}(E_2(x))$ R $\widehat{K}(E_2(y))$ (similarly).

For E we have

$$K(x) \ E \ \overline{E_2}(K(y)) \leftrightarrow K(x) \ \overline{E_2} \ K(y) \leftrightarrow x \ E_2 \ y \leftrightarrow x \ E \ E_2(y) \leftrightarrow \widehat{K}(x) \ E \ \widehat{K}(E_2(y)).$$

Thus, f is an isomorphism of MS4-frames.

LEMMA 6.13. Let $g_1, ..., g_n \in T(\mathfrak{F})^*$ and L be the correct partition of $T(\mathfrak{F})$ corresponding to the n-generated subalgebra $\mathfrak{B} = \langle g_1, ..., g_n \rangle_{T(\mathfrak{F})^*}$ of $T(\mathfrak{F})^*$. Let K be the correct partition of \mathfrak{F} corresponding to the 2n-generated subalgebra

$$\mathfrak{B}' = \langle g_1 \cap D_2, E(g_1 \cap D_1) \cap D_2, \dots, g_n \cap D_2, E(g_n \cap D_1) \cap D_2 \rangle_{\mathfrak{F}^*}$$

of \mathfrak{F}^* . Then $\widehat{K} \subseteq L$.

PROOF. It is sufficient to show that each generator of \mathfrak{B} is \widehat{K} -saturated. Since \widehat{K} relates no points across layers, and each $g_i \cap D_2$ is K-saturated (as a generator of \mathfrak{B}'), we need only check that $g_i \cap D_1$ is \overline{K} -saturated. Suppose $\alpha \in g_i$ and $\alpha \overline{K}\beta$. Then there are $x \in \alpha$, $y \in \beta$ with xKy. Evidently, $x \in E(g_i \cap D_1) \cap D_2$ which is K-saturated (as a generator of \mathfrak{B}'), so $y \in E(g_i \cap D_1) \cap D_2$. That is, y is E-related to something in $g_i \cap D_1$, which could only be β , and hence $\beta \in g_i$.

We now adapt our translation on frames to one on varieties of algebras. For $\mathfrak{A} \in \mathbf{S5}_2$, let

$$T(\mathfrak{A}) = (T(\mathfrak{A}_*))^*.$$

We write H, S, and P for the class operators of taking homomorphic images, subalgebras, and products, respectively. For a variety $V \subseteq S5_2$, let V_{SI} be the class of s.i. V-algebras, and define T(V) to be the variety generated by the translations of the s.i. members:

$$T(\mathbf{V}) = HSP(\{T(\mathfrak{A}) : \mathfrak{A} \in \mathbf{V}_{SI}\}).$$

It follows from Lemma 6.10 that T(V) is generated by a class of MS4_S[2]-algebras, so $T(V) \subseteq MS4_S[2]$.

THEOREM 6.14. T preserves and reflects local finiteness; that is, $V \subseteq S5_2$ is locally finite iff $T(V) \subseteq MS4_S[2]$ is locally finite.

PROOF. Suppose $T(\mathbf{V})$ is locally finite, and let $\mathfrak{A} = \mathfrak{F}^* \in \mathbf{V}_{SI}$ be generated by g_1, \ldots, g_n . Then each g_i is a subset of $D_2 \subseteq T(\mathfrak{F})$, and by Lemma 6.11 the operations of $\mathfrak{A}_2 \cong \mathfrak{A}$ are definable using the clopen set D_2 . By assumption, there is $f: \omega \to \omega$ such that the size of the subalgebra $\langle g_1, \ldots, g_n, D_2 \rangle \subseteq T(\mathfrak{A})$ is bounded by f(n+1). This subalgebra contains every element of \mathfrak{A} . Thus, $|\mathfrak{A}| \leq f(n+1)$. By [5, Theorem 3.7(4)], we conclude that \mathbf{V} is locally finite.

Suppose V is locally finite. To prove that T(V) is locally finite, it is sufficient to show that the generating class $\{T(\mathfrak{A}): \mathfrak{A} \in V_{SI}\}$ is uniformly locally finite (see [26, p. 285]). Suppose \mathfrak{B} is an n-generated subalgebra of $T(\mathfrak{A})$, corresponding to a correct partition L of $T(\mathfrak{A}_*)$. By Lemma 6.13, there is a 2n-generated subalgebra $\mathfrak{B}' \subseteq \mathfrak{A}$ such that the corresponding correct partition K of \mathfrak{A}_* satisfies $\widehat{K} \subseteq L$. Therefore, $\mathfrak{B} \subseteq T(\mathfrak{B}')$. By assumption, we have $f: \omega \to \omega$ such that $|\mathfrak{B}'| \le f(2n)$. Thus, it follows from Construction 6.8 that the size of $T(\mathfrak{B}')$ and hence of \mathfrak{B} is bounded by $f(2n)^2$.

§7. Conclusion and future work. We have shown that, though the Segerberg–Maksimova theorem generalizes to certain subvarieties of MS4, it fails drastically for MS4 in general, and already in MS4_S. As shown in Section 6, characterizing local finiteness even in MS4_S[2] is at least as hard as in S5₂. It is then natural to ask if the problem of local finiteness in MS4 is *strictly* harder; i.e., would a characterization of local finiteness in S5₂ yield one for MS4? We finish with an example that suggests a negative answer.

The frame $T(\mathfrak{F})$ constructed in Section 6 has the property that the relativization \mathfrak{A}_1 to D_1 is a locally finite S5²-algebra, while the relativization \mathfrak{A}_2 to D_2 is the original S5₂-algebra \mathfrak{F}^* . In light of Theorem 6.14, the algebra $T(\mathfrak{F})^*$ is locally finite iff \mathfrak{A}_2 is locally finite, which happens iff the relativization to each of the layers of $T(\mathfrak{F})$ is locally finite. It is then natural to conjecture that the local finiteness of an MS4-algebra (of finite depth) is completely determined by the local finiteness of the S5₂-algebras of the relativizations to each layer. We show this is not the case already in MS4_S[3]. Indeed, the following figure gives an example of an MS4_S[3]-frame \mathfrak{F} whose relativization to each of the three layers is a locally finite S5₂-algebra, and yet \mathfrak{F}^* fails to be locally finite:

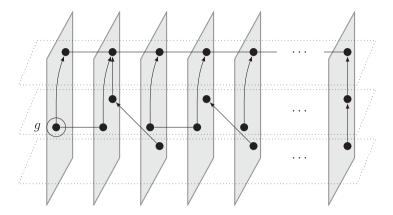


FIGURE 6. An MS4_S[3]-frame \mathfrak{F} in which the relativization to each layer is a locally finite S5₂-algebra, but \mathfrak{F}^* is not locally finite.

EXAMPLE 7.1. Let \mathfrak{F} be the frame in Figure 6. Topologically, \mathfrak{F} is the disjoint union of three layers, where the topology on each layer makes the right-most point the limit point of a one-point compactification of the finite-index points with the discrete topology. An argument similar to the proof of Lemma 6.10 shows that R and E are continuous relations, and the "top rail" ensures that \mathfrak{F} is an MS4-frame. Clearly the R-depth of \mathfrak{F} is 3.

Let g be the singleton set depicted in the figure and d the top layer. Then we can generate an infinite family of sets from the two clopen sets g and d as follows:

$$s_0 = g$$
 $s_{n+1} = \exists \Diamond s_n - d.$

Thus, \mathfrak{F}^* is not locally finite.

However, the relativization \mathfrak{A}_i to each layer is easily seen to be a locally finite S5₂-algebra. As remarked in the discussion preceding Lemma 6.11, \mathfrak{A}_1 belongs to a proper subvariety of S5², hence is locally finite. The second and third layers, equipped with the restrictions of the relations R and E, form an S5₂-frame that is a compactification of an infinite disjoint union of a single finite frame (a 3-element frame in the case of the second layer, and a singleton frame in the case of the third). Thus, for $i \in \{2, 3\}$, \mathfrak{A}_i is a subalgebra of an infinite power of a single finite algebra, so \mathfrak{A}_i belongs to a finitely generated variety, and hence \mathfrak{A}_i is locally finite.

There are two natural directions of research suggested by the results in this paper. First, as was demonstrated in [2], the monadic version of Casari's formula plays an important role in obtaining a faithful provability interpretation of the one-variable fragment of intuitionistic predicate logic. Let M^+S4 be the extension of MS4 obtained by postulating the Gödel translation of this formula. The restrictions on the interaction of the relations R and E in M^+S4 -frames make it plausible to expect a more manageable characterization of locally finite varieties of M^+S4 -algebras akin the Segerberg–Maksimova theorem.

Another direction of research is suggested by the nature of the non-locally finite frames given in the paper. The S5₂-frame in Figure 4 and the frame in Figure 6 both exhibit a one-generated infinite subalgebra in a similar manner, by alternating through the (non-commuting) relations. It is worth investigating whether this phenomenon is the only reason for non-local finiteness in S5₂, and even in MS4.

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